

Lecture 1
EVOLUTIONARY GAME THEORY
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1 Economic theory and "as if" rationality

- The rationalistic paradigm in economics: Savage rationality [Leonard Savage: *The Foundations of Statistics*, 1954]
 - Each economic agent's behavior derived from maximization of some goal function (utility, profit), under given constraints and information
- The "as if" defence by Milton Friedman (1953): *The methodology of positive economics*
 - Firms that do not take profit-maximizing actions are selected against in the market
 - But is this claim right? Under perfect competition? Under imperfect competition?

- Evolutionary theorizing: De Mandeville, Malthus, Darwin, Maynard Smith
- Darwin: exogenous environment - “perfect competition”
- Maynard Smith: endogenous environment - “imperfect competition”

2 Game theory

- *A mathematically formalized theory of strategic interaction.*

- Applications abound,

— in economics

— in political science

— in biology

— in computer science

3 John Nash's two interpretations of equilibrium

Nash (Ph.D. thesis, Mathematics Department, Princeton, 1950):

1. The rationalistic (or epistemic) interpretation
2. The "mass action" (or evolutionary) interpretation

3.1 The rationalistic interpretation

1. The players interact exactly once (true even if the interaction itself is a repeated game).
2. The players are *rational* in the sense of Savage (1954)
3. Each player *knows* the game in question (strategy sets, payoffs to all players)

However, this clearly does not imply that they will play a Nash equilibrium. Hence, assume in addition that

4. The game and all players' rationality is *common knowledge*

Example 3.1 (Coordination game)

	<i>L</i>	<i>R</i>
<i>L</i>	2, 2	0, 0
<i>R</i>	0, 0	1, 1

Example 3.2 (Matching Pennies)

	<i>H</i>	<i>T</i>
<i>H</i>	1, -1	-1, 1
<i>T</i>	-1, 1	1, -1

Example 3.3 *A game with a unique Nash equilibrium. This equilibrium is strict; any deviation is costly.*

	<i>L</i>	<i>C</i>	<i>R</i>
<i>T</i>	7, 0	2, 5	0, 7
<i>M</i>	5, 2	3, 3	5, 2
<i>B</i>	0, 7	2, 5	7, 0

Conclusion

The rationalistic interpretation, even under CK of the game and rationality, does not in general imply equilibrium play

[In fact, this combined hypothesis implies rationalizability.]

3.2 Nash's mass-action interpretation

1. For each *player role* in the game: a large population of identical individuals
2. The game is recurrently played, at times $t = 0, 1, 2, 3, \dots$ (or at Poisson arrival times) by randomly drawn individuals, one from each *player population*
3. Individuals play *pure strategies*
4. Individuals observe *empirical samples* of earlier behaviors and outcomes and *avoid suboptimal actions*

[According to Nash (1950), they need not know the strategy sets or payoff functions of other player roles. They do not even have to know they play a game.]

- A population-statistical distribution over the pure strategies in a player role constitute a mixed strategy
- Nash's (informal) claim: If all individuals avoid suboptimal pure strategies, and the population distribution is stationary ("stable"), then it constitutes a Nash equilibrium
 - Statistical independence across player populations \Rightarrow joint probability distribution over pure-strategy profiles constitutes a mixed-strategy profile
 - Zero probability for suboptimal pure strategies \Rightarrow each player-population's mixed strategy is a best reply
- Reconsider the above examples in this (informal) interpretation

Conclusion

The mass-action interpretation does not in general imply equilibrium play, but it almost does it, and holds promise

- In fact, *evolutionary game theory* provides methods and concepts to rigorously explore these promises
- The “folk theorem” of evolutionary game theory:
 - If the population process *converges* from an interior initial state, then the limit distribution is a Nash equilibrium
 - If a stationary population distribution is *stable*, then it constitutes a Nash equilibrium

- If the population process is "*convex-monotone*," then non-rationalizable strategies have zero asymptotic probability, in all finite two-player games

4 Evolutionary game theory

- Evolutionary process =
= mutation process + selection process
 - The unit of selection: usually strategies ("strategy evolution"), sometimes utility functions ("preference evolution")
1. **Evolutionary stability:** focus on robustness to mutations
 2. **Replicator dynamic:** focus on selection. [Robustness to mutations by way of dynamic stability]
 3. **Stochastic stability:** both selection and mutations

5 Evolutionary stability under strategy evolution

- ESS = *evolutionarily stable strategy* [Maynard Smith and Price (1972), Maynard Smith (1973)]
 - “a strategy that ‘cannot be overturned’, once it has become the ‘convention’ in a population

Heuristically

1. A large population of individuals who are recurrently and (uniformly) randomly matched in pairs to play a finite and symmetric two-player game
2. Initially, all individuals always use the same pure or mixed strategy, x , the *incumbent* (or *resident*) strategy
3. Suddenly, a small population share $\varepsilon > 0$ switch to another pure or mixed strategy, y , the *mutant* strategy
4. If the incumbents (residents) on average *do better* than the mutants, then x is *evolutionarily stable against* y

5. x is called *evolutionarily stable* if it is evolutionarily stable against *all* mutations $y \neq x$

Domain of analysis

- *Symmetric finite two-player games in normal form*

Definition 5.1 *A finite and symmetric two-player game is any normal-form game $G = (N, S, u)$ with $N = \{1, 2\}$, $S_1 = S_2 = S = \{1, \dots, m\}$ and $u_2(h, k) = u_1(k, h)$ for all $h, k \in S$.*

- Payoff bimatrix (A, B) , where $A = (a_{hk})$, $B = (b_{hk})$
- Game symmetric iff $B = A^T$.

Example 5.1 (Prisoners' dilemma) *Payoff bimatrix:*

	<i>C</i>	<i>D</i>
<i>C</i>	3, 3	0, 4
<i>D</i>	4, 0	2, 2

$$A = \begin{pmatrix} 3 & 0 \\ 4 & 2 \end{pmatrix} \quad B = \begin{pmatrix} 3 & 4 \\ 0 & 2 \end{pmatrix}$$

Symmetric since $B = A^T$.

- Write Δ for $\Delta(S)$, the mixed-strategy simplex:

$$\Delta = \{x \in \mathbb{R}_+^m : \sum_{i \in S} x_i = 1\}$$

- Write the payoff to any strategy $x \in \Delta$, when used against any strategy $y \in \Delta$ as

$$\pi(x, y) = x^T A y$$

Note that the first argument is own strategy, second argument the other party's strategy.

- While the prisoner dilemma is symmetric, matching-pennies is not

Example 5.2 (Matching Pennies) *Payoff bimatrix:*

	<i>H</i>	<i>T</i>
<i>H</i>	1, -1	-1, 1
<i>T</i>	-1, 1	1, -1

$$A = \begin{pmatrix} 1 & -1 \\ -1 & 1 \end{pmatrix} \quad B = \begin{pmatrix} -1 & 1 \\ 1 & -1 \end{pmatrix}$$

Here $B^T \neq A$. Not a symmetric game.

- Thus, matching pennies games fall outside the domain of evolutionary stability analysis.
- However, if player roles are randomly assigned, with equal probability for both role-allocations, then the so-defined metagame is symmetric, and evolutionary stability analysis applies to the metagame.

Example 5.3 (Coordination game) *Payoff bimatrix:*

	<i>L</i>	<i>R</i>
<i>L</i>	2, 2	0, 0
<i>R</i>	0, 0	1, 1

$$A = B = \begin{pmatrix} 2 & 0 \\ 0 & 1 \end{pmatrix}$$

A doubly symmetric game: $B = A^T = A$ (this is an example of a 'potential game', Monderer and **Shapley**, 1996)

- Best replies to $x \in \Delta$:

$$\beta^*(x) = \{x^* \in \Delta : \pi(x^*, x) \geq \pi(x', x) \quad \forall x' \in \Delta\}$$

- This defines a *correspondence* from Δ to itself: $\beta^* : \Delta \rightrightarrows \Delta$

[This is distinct from the usual BR correspondence, which maps Δ^2 to Δ]

- Let

$$\Delta^{SNE} = \{x \in \Delta : x \in \beta^*(x)\}$$

- Note $x \in \Delta^{SNE} \Leftrightarrow (x, x)$ is a symmetric Nash equilibrium

Proposition 5.1 $\Delta^{SNE} \neq \emptyset$.

Proof: Application of Kakutani's Fixed-Point Theorem.

5.1 Definition of ESS

- We are now in a position to define evolutionary stability exactly:

Definition 5.2 $x \in \Delta$ is an evolutionarily stable strategy (ESS) if for every strategy $y \neq x \exists \bar{\varepsilon}_y \in (0, 1)$ such that for all $\varepsilon \in (0, \bar{\varepsilon}_y)$:

$$\pi [x, \varepsilon y + (1 - \varepsilon)x] > \pi [y, \varepsilon y + (1 - \varepsilon)x].$$

- “Post-entry mixture”:

$$p = \varepsilon y + (1 - \varepsilon)x \in \Delta$$

a convex combination of x and y , a point on the straight line between them.

- Let $\Delta^{ESS} \subset \Delta$ denote the set of ESSs
- Note that an ESS has to be a *best* reply to itself: if $x \in \Delta^{ESS}$ then $\pi(y, x) \leq \pi(x, x)$ for all $y \in \Delta$
- Hence $\Delta^{ESS} \subset \Delta^{SNE}$
- Note also that an ESS has to be a *better* reply to its alternative best replies than they are to themselves: if $x \in \Delta^{ESS}$, $y \in \beta^*(x)$ and $y \neq x$, then $\pi(x, y) > \pi(y, y)$

Proposition 5.2 $x \in \Delta^{ESS}$ if and only if for all $y \neq x$:

$$\pi(y, x) \leq \pi(x, x)$$

and

$$\pi(y, x) = \pi(x, x) \Rightarrow \pi(y, y) < \pi(x, y)$$

5.2 Examples

5.2.1 Prisoner's dilemma

	<i>C</i>	<i>D</i>
<i>C</i>	3, 3	0, 4
<i>D</i>	4, 0	2, 2

$$\Delta^{ESS} = \Delta^{SNE} = \{D\}$$

5.2.2 Coordination game

	<i>L</i>	<i>R</i>
<i>L</i>	2, 2	0, 0
<i>R</i>	0, 0	1, 1

$$\Delta^{SNE} = \left\{ L, R, \frac{1}{3}L + \frac{2}{3}R \right\}$$

$$\Delta^{ESS} = \{L, R\}$$

The mixed NE is *perfect* and even *proper*, but not evolutionarily stable

5.2.3 Partnership game

- Small start-up businesses with two partners, or pairs of students assigned to write an essay together
- Each partner has to choose between *work* (or “contribute”) and *shirk* (or “free-ride”)
 - If both choose *W*: *good outcome for both*
 - If one chooses *W* and the other *S*: *loss to the first but gain to the second*
 - If both choose *S*: *heavy loss to both*
- Example of payoff bimatrix:

	<i>W</i>	<i>S</i>
<i>W</i>	3, 3	0, 4
<i>S</i>	4, 0	-1, -1

- Symmetric game, but **not** a Prisoners' Dilemma: S does not dominate W: it is better to work if the other shirks
- Consider a large pool of individuals and random matching
- What do you think would happen? A tendency to work? To shirk? Any tendency at all?
- Any ESS? This is strategically equivalent to a so-called *hawk-dove game*, where work = dove and shirk = hawk

1. Unique symmetric NE: randomize uniformly, $x^* = (1/2, 1/2)$, $\Delta^{SNE} = \{x^*\}$. Hence $\Delta^{ESS} \subset \{x^*\}$

2. x^* an ESS iff

$$\pi(x^*, y) > \pi(y, y) \quad \forall y \neq x^*$$

3. Equivalently:

$$\frac{1}{2} [3y_1 + 4y_1 - (1 - y_1)] > 3y_1^2 + 4y_1(1 - y_1) - (1 - y_1)^2$$

or

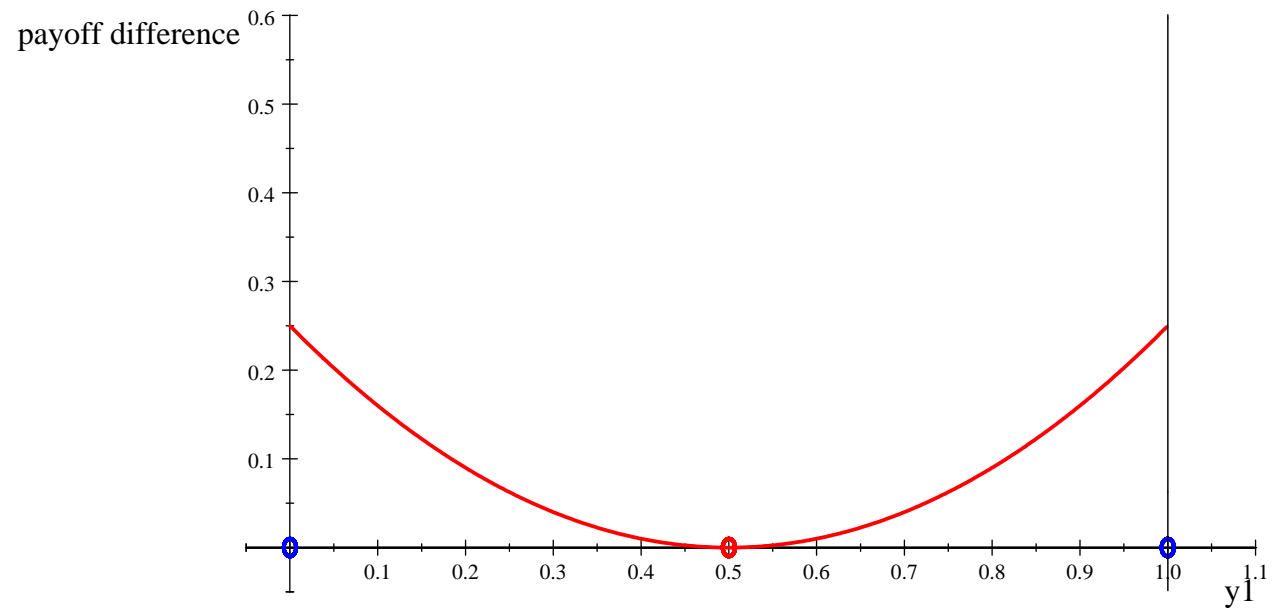
$$8y_1 - 1 > -4y_1^2 + 12y_1 - 2$$

or

$$4 \left(y_1 - \frac{1}{2} \right)^2 > 0$$

4. True, hence x^* is an ESS!

- Graphical illustration of payoff difference $\pi(x^*, y) - \pi(y, y)$:



- Some games have no ESS. For instance, when all payoffs are the same. But also in more interesting games such as

Example 5.4 (Rock-scissors-paper) *Rock beats Scissors, Scissors beat Paper, and Paper beats Rock:*

$$A = \begin{pmatrix} 0 & 1 & -1 \\ -1 & 0 & 1 \\ 1 & -1 & 0 \end{pmatrix}$$

Unique Nash equilibrium: $x^ = \left(\frac{1}{3}, \frac{1}{3}, \frac{1}{3}\right)$. All pure strategies are best replies and do as well against themselves as x^* does against them. $\Delta^{ESS} = \emptyset$.*

THE END

Literature: Chapter 9 in van Damme (1991) or Chapter 2 in Weibull (1995).